

- course on Quantum Computers and not on Quantum Computing.
- A few words on QC. (of states)
- Quantum dynamics involves 2^n different amplitudes calculating the evolution of which is hard. [assume 2×2 ; $n=80$ and 8 bit numerical resolution per amplitude.]

$$\Rightarrow 2^{80+8} = 2^{88} \text{ bits} \quad 1 \text{ Tbyte} \sim 10^{12} \times 8 \approx 2^{40}$$

10^{10} Tbyte. just for memory at each moment in time. Not to mention # of operations.

- More formally: complexity classes P vs. NP.

P = # of steps is polynomial in input.

NP = "Nondeterministic polynomial". Given the right solution, it can be checked in polynomial time.

$$\Rightarrow P \subseteq NP$$

$P = NP$? $\Rightarrow$ open question (1 of 7 million prize problems.)

NP-complete - set of all problems that are "as hard" as any other problem in NP.

i.e. any other problem in NP, can be reduced to it.

in polynomial time.

$\Rightarrow$ if you solve an NP complete problem in polynomial time $\Rightarrow$ NP $\subseteq$ P

- (Nearest / neighbor) Ising model

$$H = \sum_{\langle i,j \rangle} J_{ij} \sigma_i \sigma_j$$

for $J_{ij} \in$ positive / negative $\Rightarrow$ ground state ferromagnetic / antiferromagnetic

J_{ij} = changes sign randomly $\Rightarrow$ spin glass

finding ss for - nearest neighbors in 3D

J_{ij} long range $\rightarrow$ connectivity,

finding the ground-state is NP complete

- Feynman (81) + maybe if we use a quantum computer, simulating quantum dynamics would be easy

- Not necessarily: in the case of the Ising G.S. ^{adiabatically} following to g.s. but in principle efficient $\leftarrow$ correlation
- meas. will collapse onto FC meas. basis.
- some global properties $\rightarrow$ i.e. g.s. magnetization spin correlation functions.

Since then: several algorithms

- Rely on QFT $\Rightarrow$ Shor's algorithm for factorization $\hookrightarrow$ not NP hard!
- Data Search $\Rightarrow$ Grover's algorithm $2^n \rightarrow \sqrt{2^n}$
- Q Simulations
- Scott Aaronson: sat. matrix permanent - NP complete using Q random walk

- complexity class BQP $\hookrightarrow$ Bounded probability.

All these rely on a Quantum parallelism

+ b. Interference to get one (global) answer.

Example: Deutsch - Jozsa

$\{0,1\} \rightarrow \{0,1\}$ classically 2 queries to learn that function.

$$U_f |x\rangle |y\rangle = |x\rangle |y \oplus f(x)\rangle$$

unitary.

$$\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) |0\rangle = \frac{1}{\sqrt{2}} [|0\rangle |f(0)\rangle + |1\rangle |f(1)\rangle]$$

Get both qubits in the same superposition.

Not enough $\Rightarrow$ have to measure at least twice.

- Global question.

$$f(0) = f(1) = ?$$

$$\text{or } f(0) = \overline{f(1)}$$

$$\frac{1}{\sqrt{2}} [|0\rangle + |1\rangle] \frac{1}{\sqrt{2}} [|0\rangle - |1\rangle] \Rightarrow$$

$$\frac{1}{2} (|0\rangle |f(0)\rangle - |0\rangle |f(1)\rangle)$$

$$\frac{1}{2} [10f(0) - 10, 1 \oplus f(0) + 11f(1) - 11, 1 \oplus f(1)]$$

$$1 \oplus 0 = 1$$

$$\Rightarrow 1 \oplus x = \bar{x}$$

$$1 \oplus 1 = 0$$

$$= \frac{1}{2} [10f_0 - 10\bar{f}_0 + 11f_1 - 11\bar{f}_1]$$

$$\text{if } f_0 = f_1 = f < \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}$$

$$= \frac{1}{2} [10 > [1f - 1\bar{f}] + 11 > [1f - 1\bar{f}]] =$$

$$= \frac{1}{2} (10 > + 11 >) (1f - 1\bar{f}) = \pm \frac{1}{2} (10 > + 11 >) (10 > - 11 >)$$

$$\text{if } f_0 = f_1 = f < \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}$$

$$f < \begin{smallmatrix} 0 \\ 1 \end{smallmatrix}$$

$$= \frac{1}{2} [10f - 10\tilde{f} + 11\tilde{f} - 11f]$$

$$\frac{1}{2} [(10 - 11)f - (10 - 11)\tilde{f}] =$$

$$= \frac{1}{2} (10 - 11) (f - \tilde{f}) = \pm \frac{1}{2} (10 - 11) (10 - 11)$$

can be extended to larger x

$f(x) < \text{constant}$
balanced?

A Universal Q computer.

~~Not a task~~ - Register of qubits (spins)

"supermarket"

- initialize

- implement ANY unitary operation

- measure

- without noise interfering

A tough job:

examples: arbit. unitary on n spins $2^n \times 2^n = 4^n$

or ind. generators $\Rightarrow$ Pauli group $\underbrace{\sigma_x \sigma_y \dots \sigma_z}_{n}$

how do you generate this?

- effect of noise better looked on using states.

$$|\psi\rangle = \frac{1}{\sqrt{2}} \left[\underbrace{|00\dots 0\rangle}_{n} + e^{i\theta} |11\dots 1\rangle \right]$$

meas. 1 collapses the full superposition!

write in the basis $|+\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle)$

$$|-\rangle = \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$\varphi = 0 \quad |\psi\rangle = \frac{1}{\sqrt{2^{n-1}}} \sum |even\rangle \quad \text{i.e. even \# of 1's}$$

$$\varphi = \pi \quad |\psi\rangle = \frac{1}{\sqrt{2^{n-1}}} \sum |odd\rangle \quad \text{i.e. odd \# of 1's}$$

$\Rightarrow$ a single $|+\rangle \rightarrow |-\rangle$ bit flip completely

dephases $|\psi\rangle$.

- A universal QC is a universal state generating machine. enabling the construction of any structure in Hilbert space.

- other applications: Metrology.

More importantly, it pushes QM towards a new-exploded era.

Course content topics:

Part 1.

- Qubit
- Qubit arrays + tensor products of Hilbert spaces.
- Entanglement
- Universal gate set.
- Effect of imperfections/noise $\Rightarrow$ Quantum operations.
- QFT, QST
- Distance measures in Hilbert space.
- Quantum Error-Correction Codes.
- Fault tolerant QC.

Part 2.

- trapped ion QC

- Superconducting QC

- Photon QC.

At least for part 1: Nielsen & Chuang.

QEC: ch7 of John Preskill's Book (online)

14 2-hour meetings. (No lecture on Dec. 19th), 3 exercises. 2 exam on 30/1

Remarks:

I am giving the traditional story.

Alternatives:

e.g.

- Adiabatic QC

- Dissipative QC

- One-way QC.

- Topological QC